

# Ridesharing: Simulator, Benchmark, and Evaluation

James Jie Pan, Guoliang Li, Juntao Hu

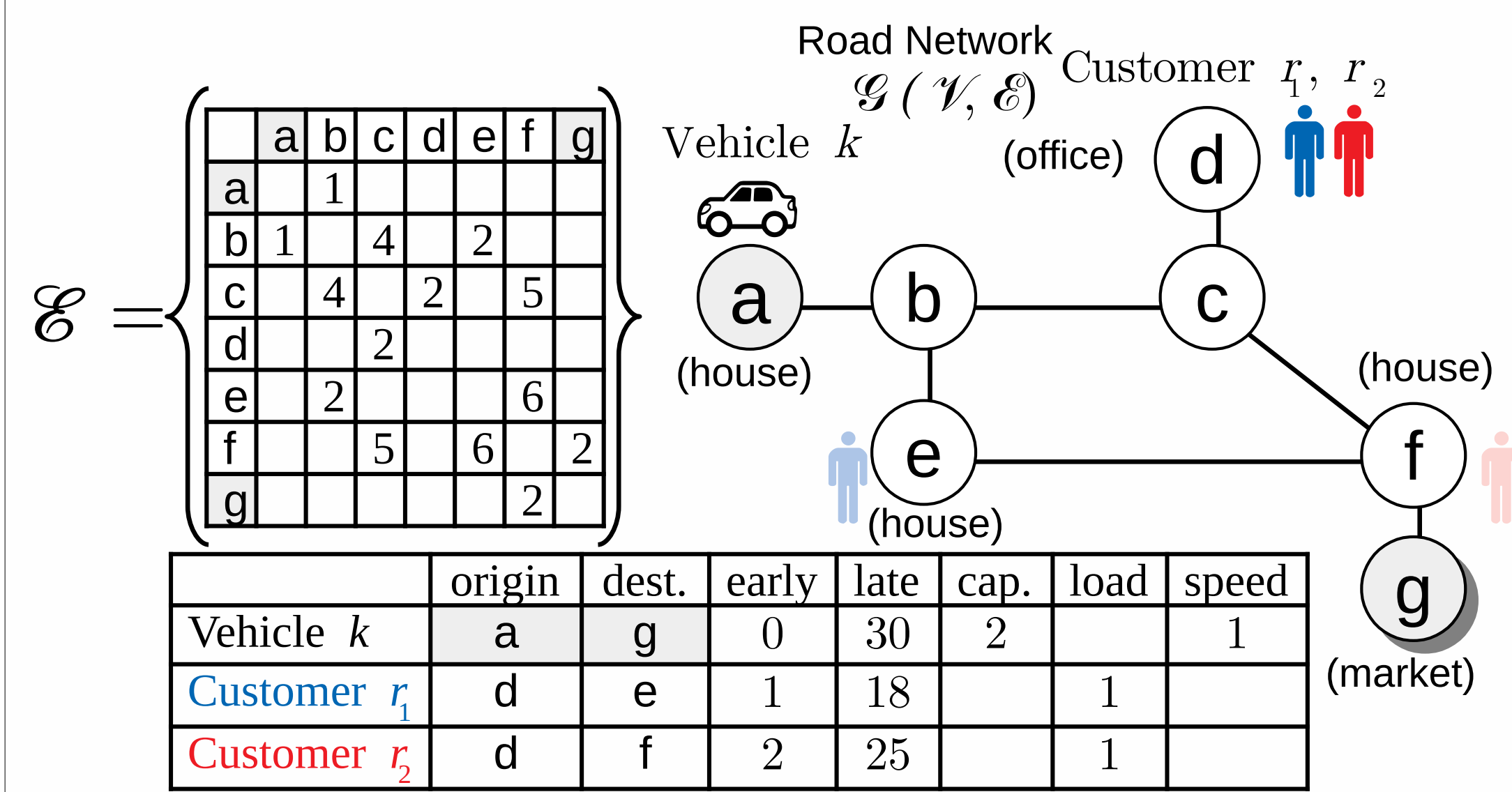
## Motivation

Can ridesharing algorithms make high-quality assignments under real-time constraint?

How can we characterize assignment techniques?

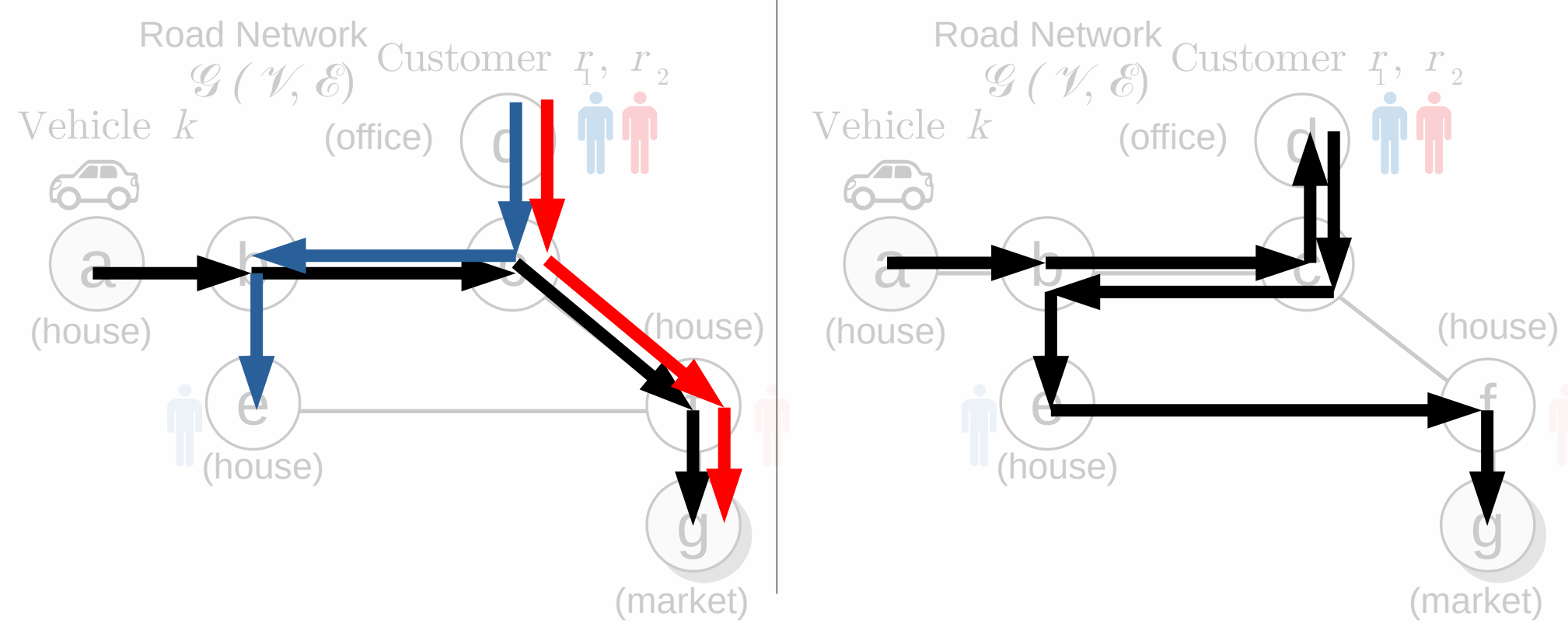
How do we design a benchmark?

## Reason for Ridesharing



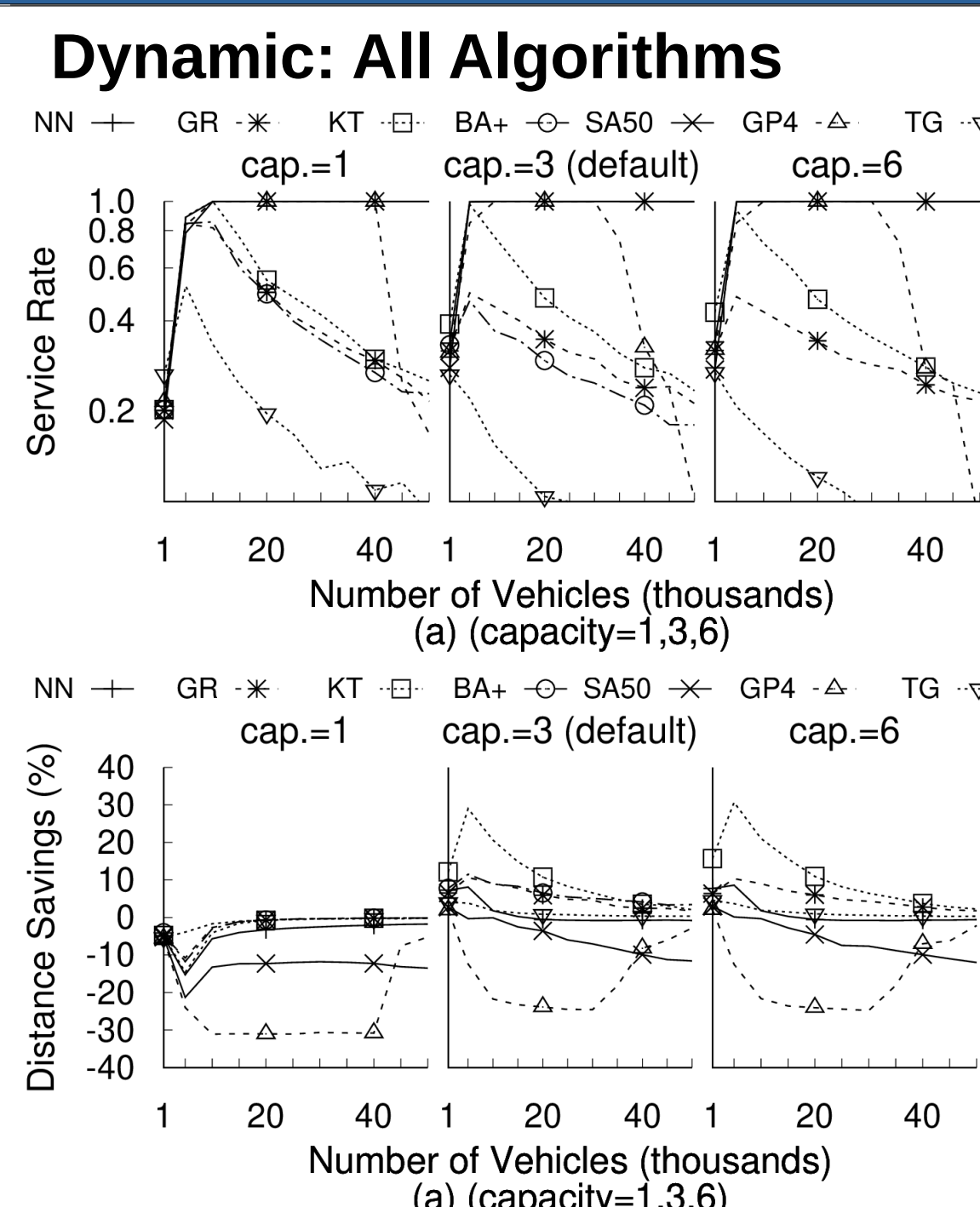
Travel Separately  
(No Ridesharing)  
Cost = 27

Travel Together  
(With Ridesharing)  
Cost = 23 ✓

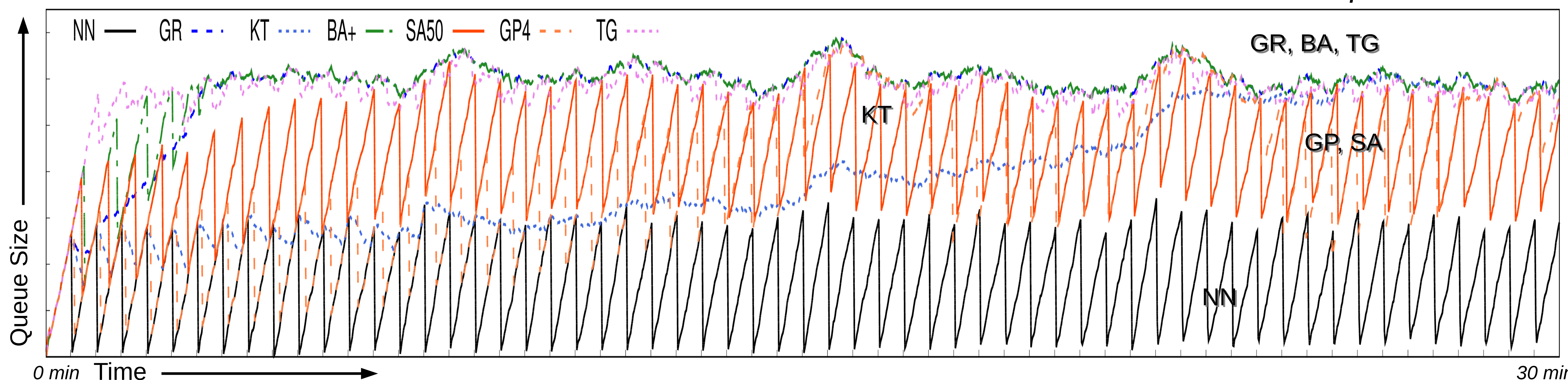


- Potential for saving total travel distances, thereby reducing overall fuel use and emissions
- Potential for reducing number of vehicles on the road

## Our Simulation Results



## Number of Queued Customers Over Time, per Algorithm



For high service rate, matching rate should meet or exceed the rate of requests.

- The algorithms that achieved the highest service rates were those that cleared the queue (Nearest Neighbor, Simulated Annealing, and GRASP).
- Kinetic Tree saved more distance, achieved better service rate, and cleared the queue faster than Greedy.
- Despite high service rate, Simulated Annealing and GRASP did not perform enough descents to result in distance savings.
- Trip-Vehicle Grouping struggled to clear the queue and had the lowest service rate.

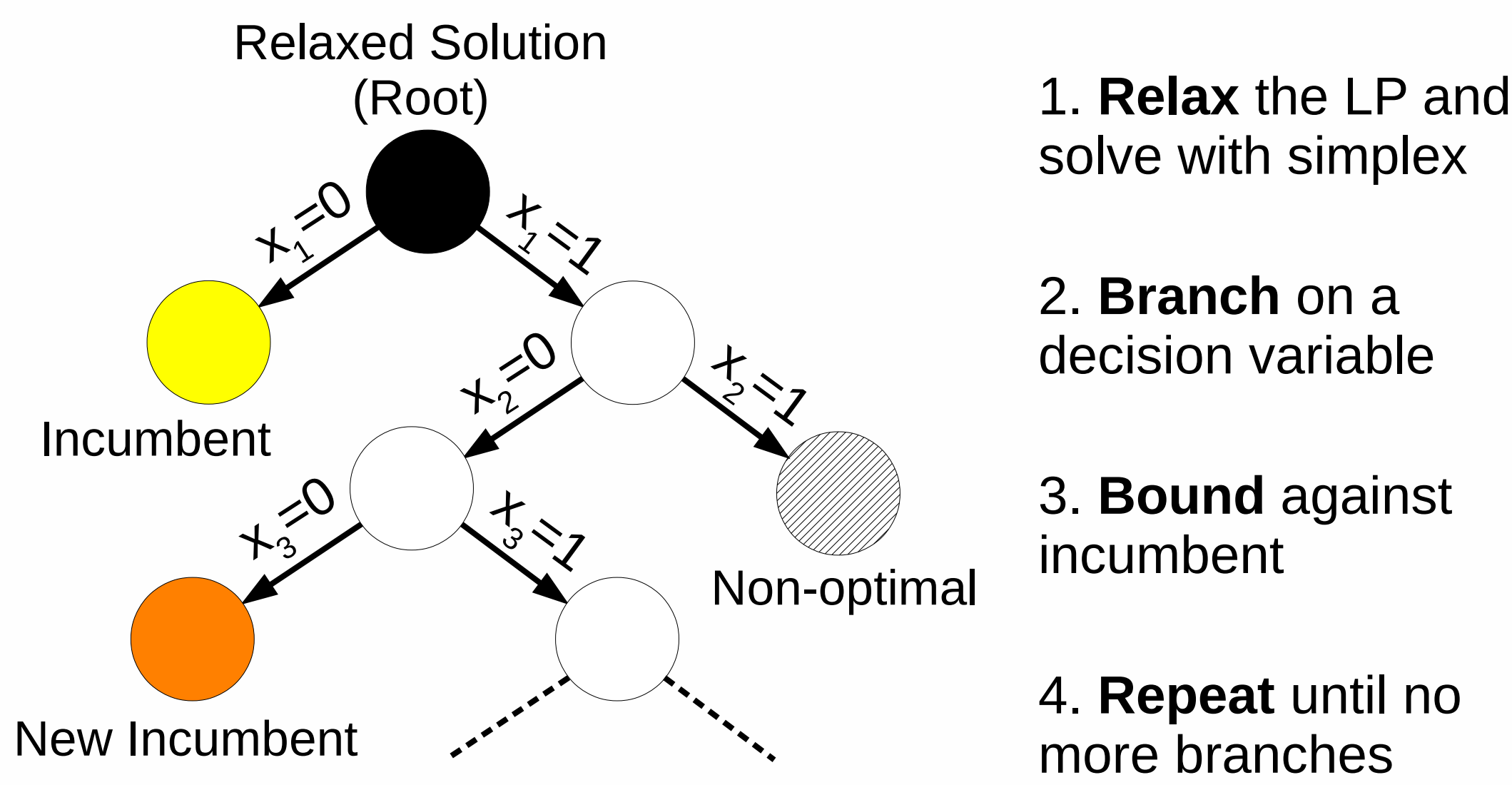
## Customer-to-Vehicle Assignment Algorithms

An Exact Offline Method: Branch-and-Bound (BB) [Cor06]

1. Label vehicles  $M = \{1..m\}$ , customers  $N = \{1..n\}$
2. Design origin/destination labels:

For...	Origin is labeled...	Destination is labeled...
Vehicle $k \in M$	$k$	$k+m$
Customer $r \in N$	$2m+r$	$2m+n+r$

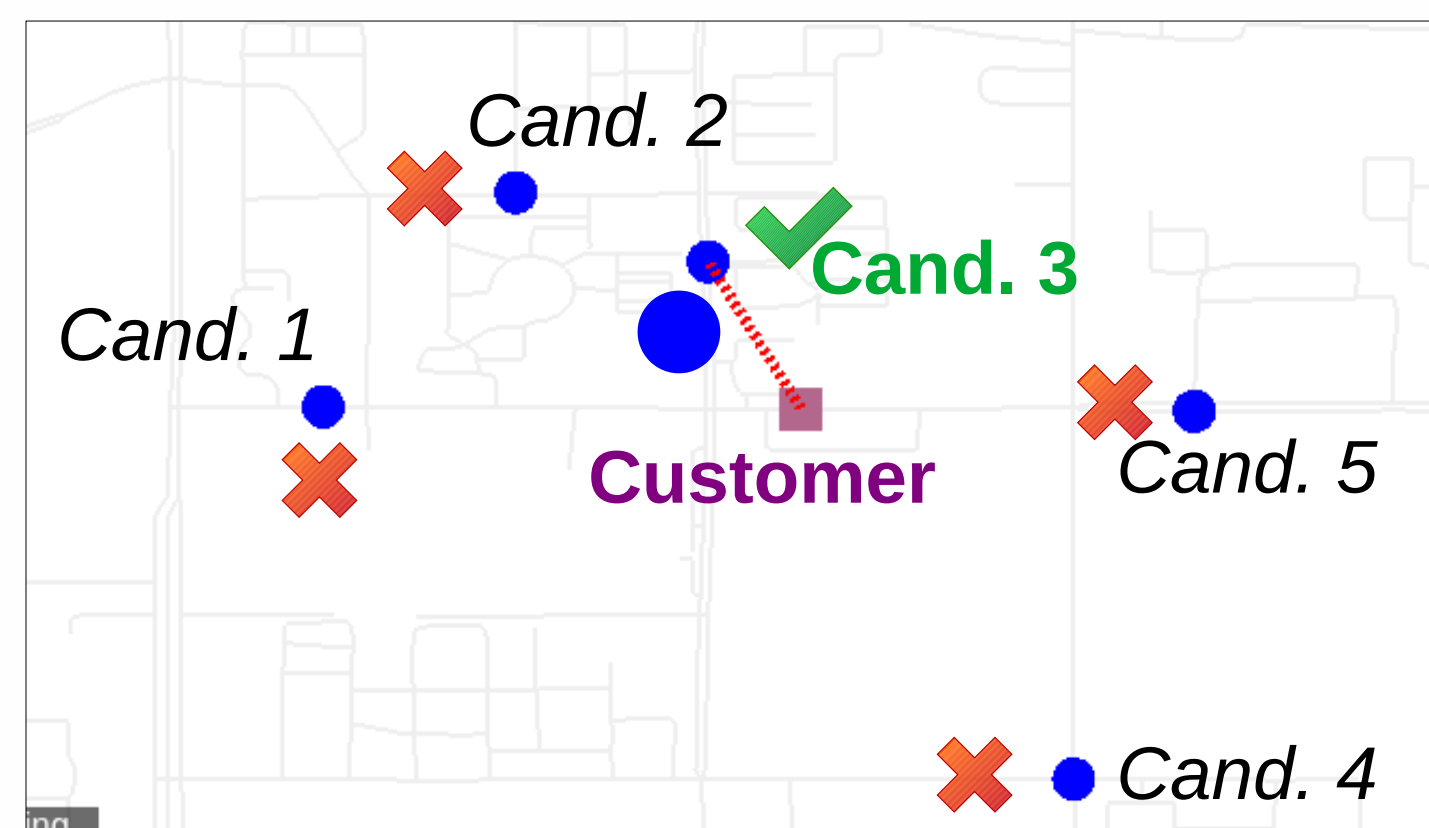
3. Min.  $\sum_{k \in M} \sum_{i \in V} \sum_{j \in V} c_{ij} x_{ij}^k$  s.t. constraints, where  $V = N_o \cup N_d \cup \{k, k+m\}$  and  $c_{ij}$  is shortest-path cost from  $i$  to  $j$ .



1. Relax the LP and solve with simplex
2. Branch on a decision variable
3. Bound against incumbent
4. Repeat until no more branches

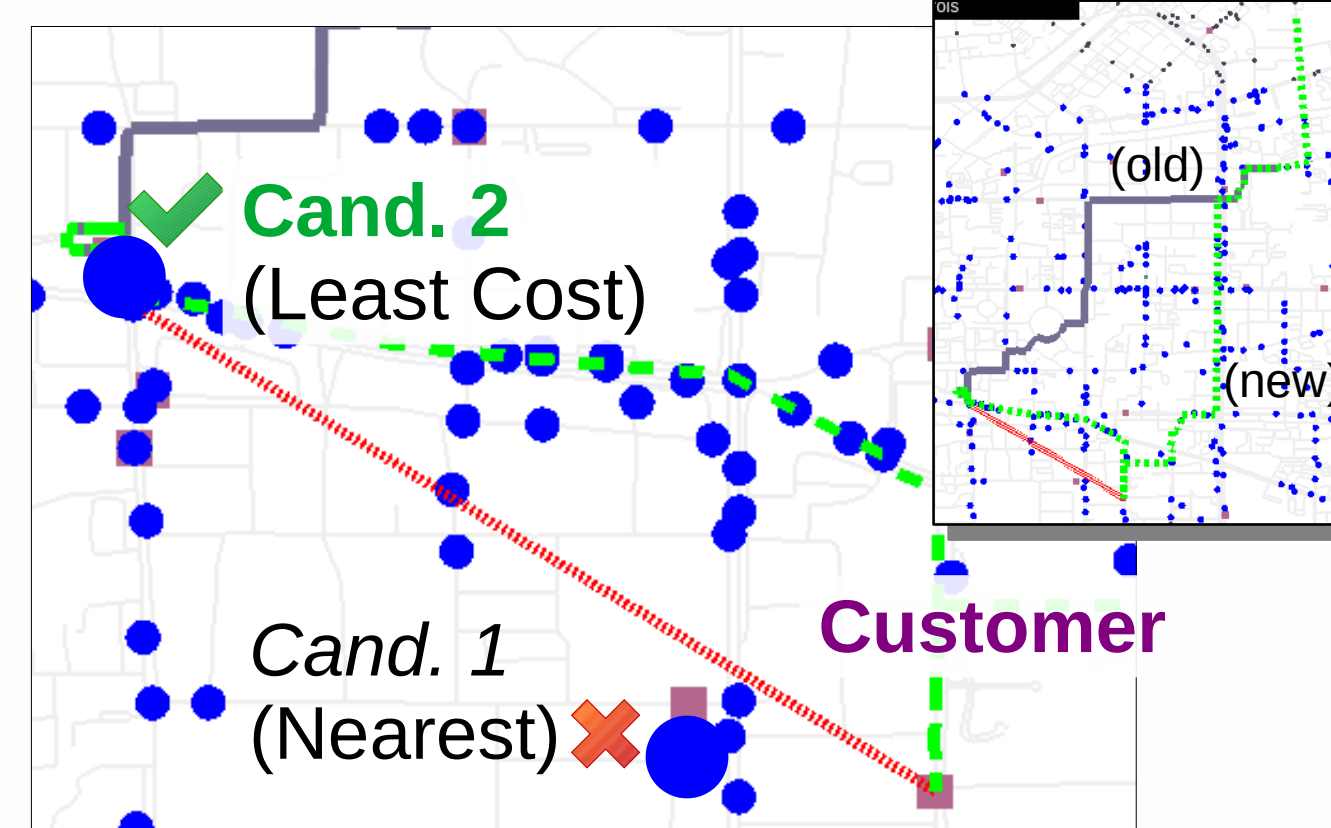
Online Search-Based Algorithms: Assignment =  $\sigma_{Pred}$ (Candidates)

Nearest-Neighbor (NN)



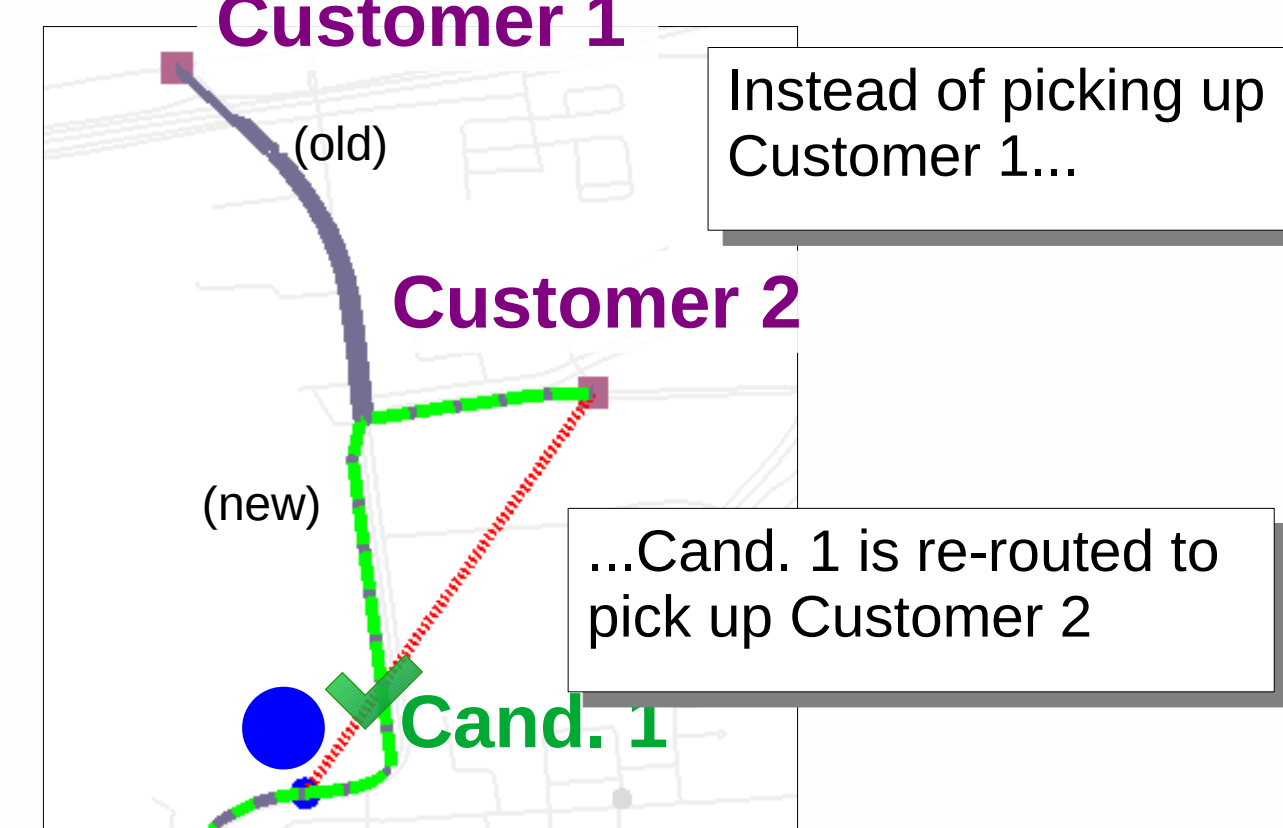
For  $k \in \text{Candidates}$ ,  
 $Pred(k) = (k \text{ is Nearest}) \wedge (k \text{ is valid})$

Greedy (GR), Kinetic Tree (KT)



For  $k \in \text{Candidates}$ ,  
 $Pred(k) = (k \text{ minimizes Cost}) \wedge (k \text{ is valid})$

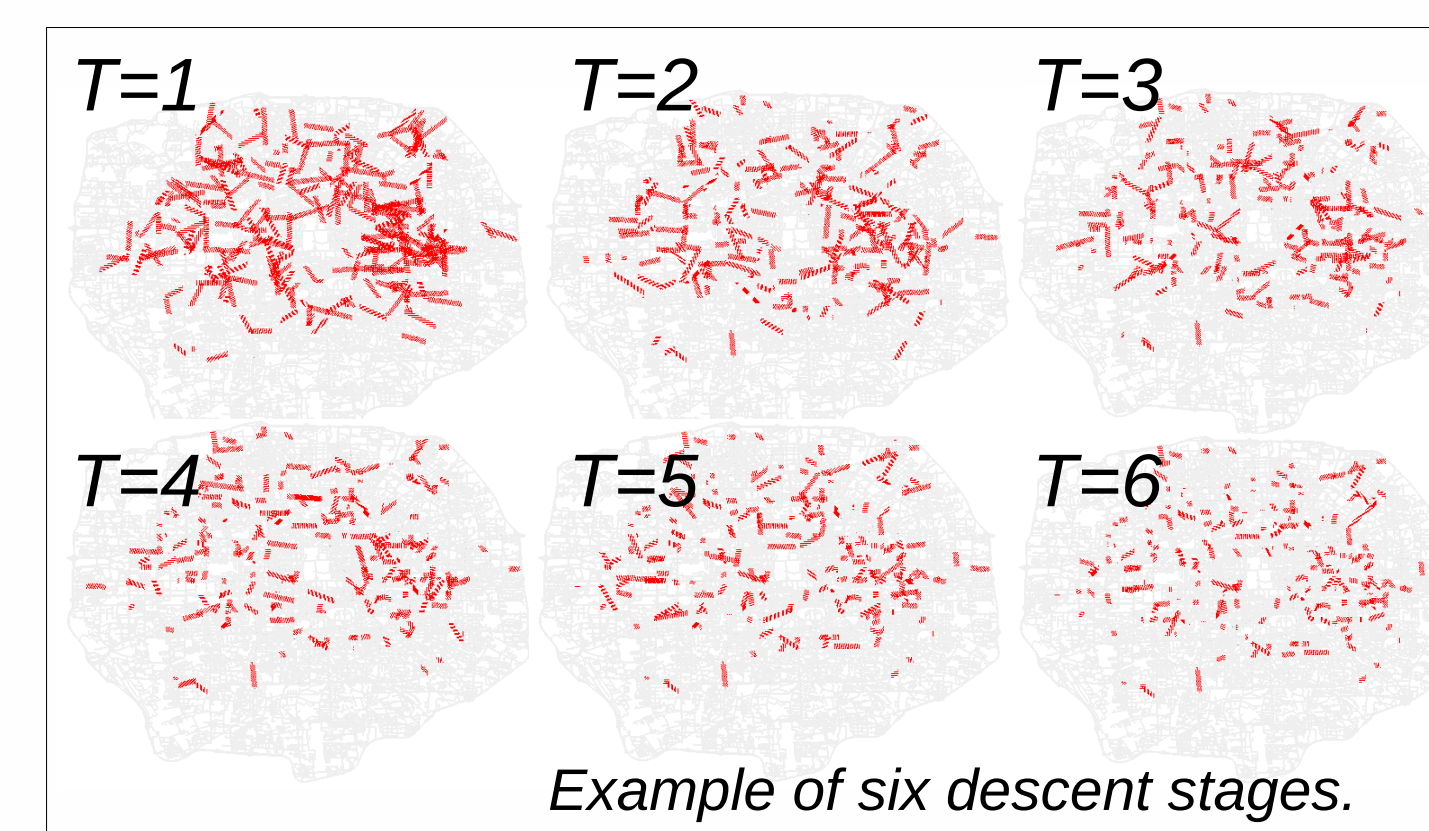
Bilateral Arrangement (BA)



For  $k \in \text{Candidates}$ ,  
 $Pred(k) = (k \text{ minimizes Cost}) \wedge (k \text{ is valid, with Replace if necessary})$

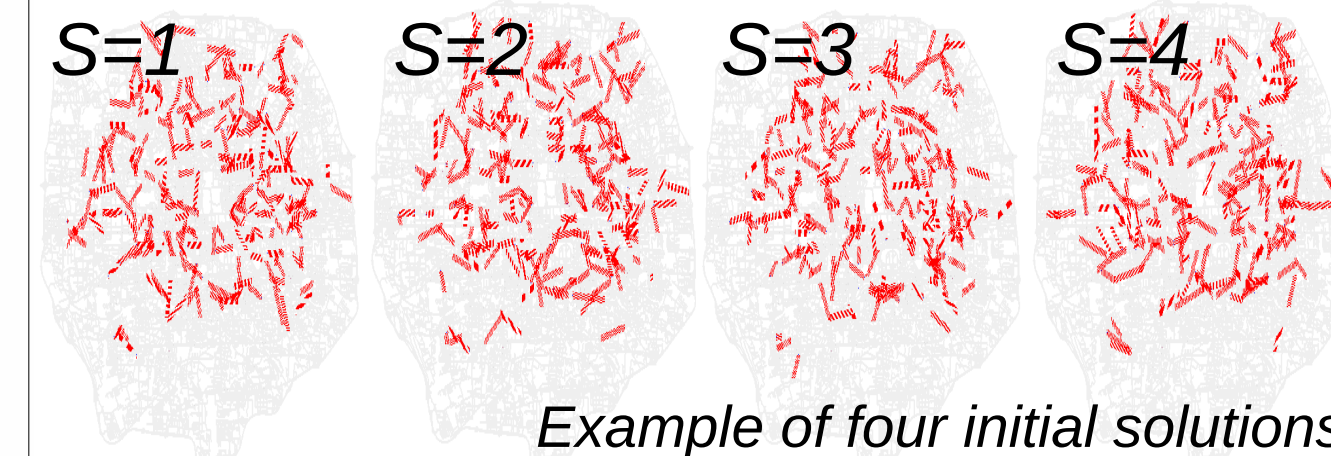
Online Join-Based Algorithms: Assignments = Customers  $\bowtie_{Pred}$  Candidates

Simulated Annealing (SA)



Descend using random reassigns, with chance of hill-climbing

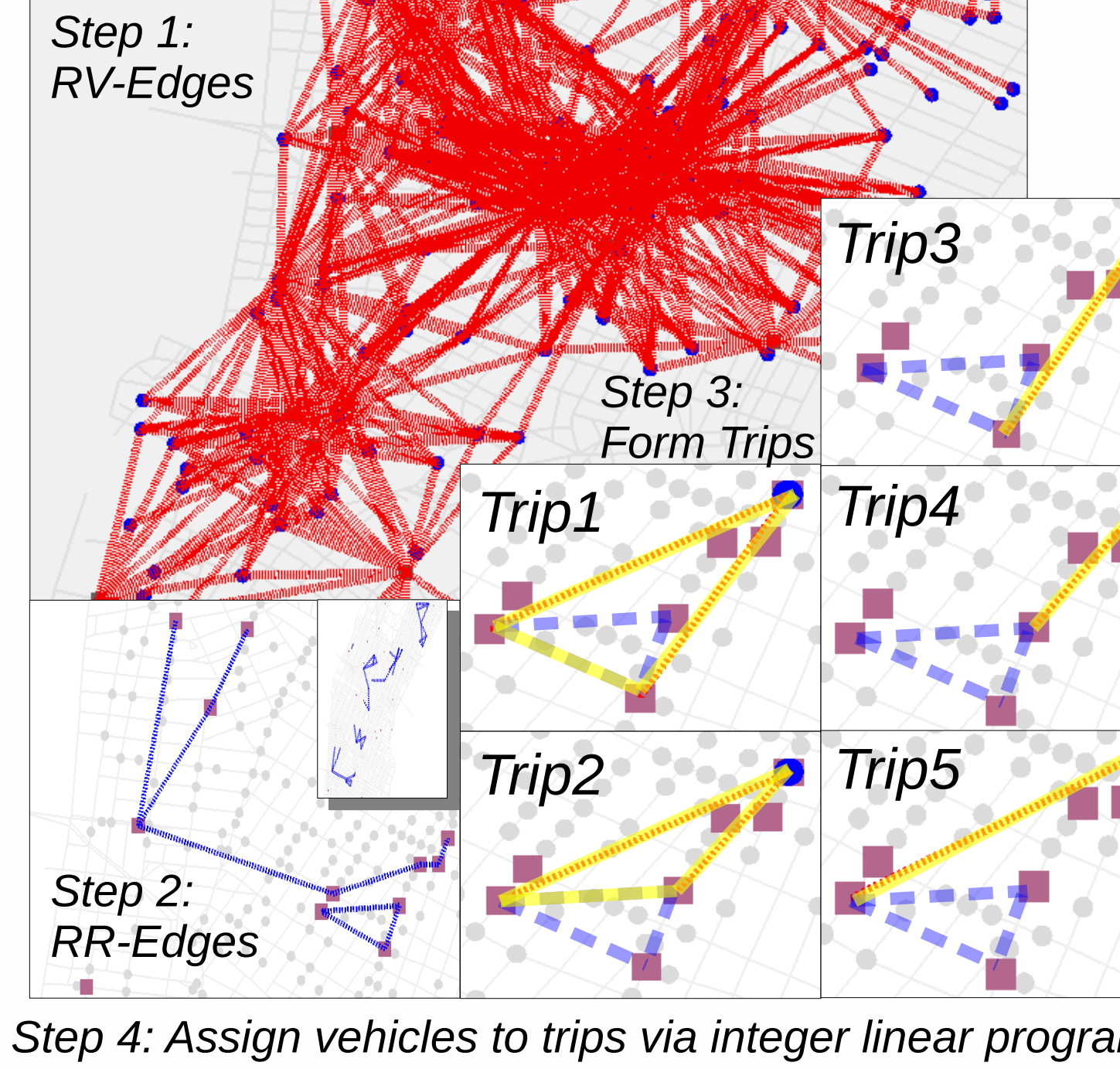
Greedy Randomized Adaptive Search Procedure (GRASP)



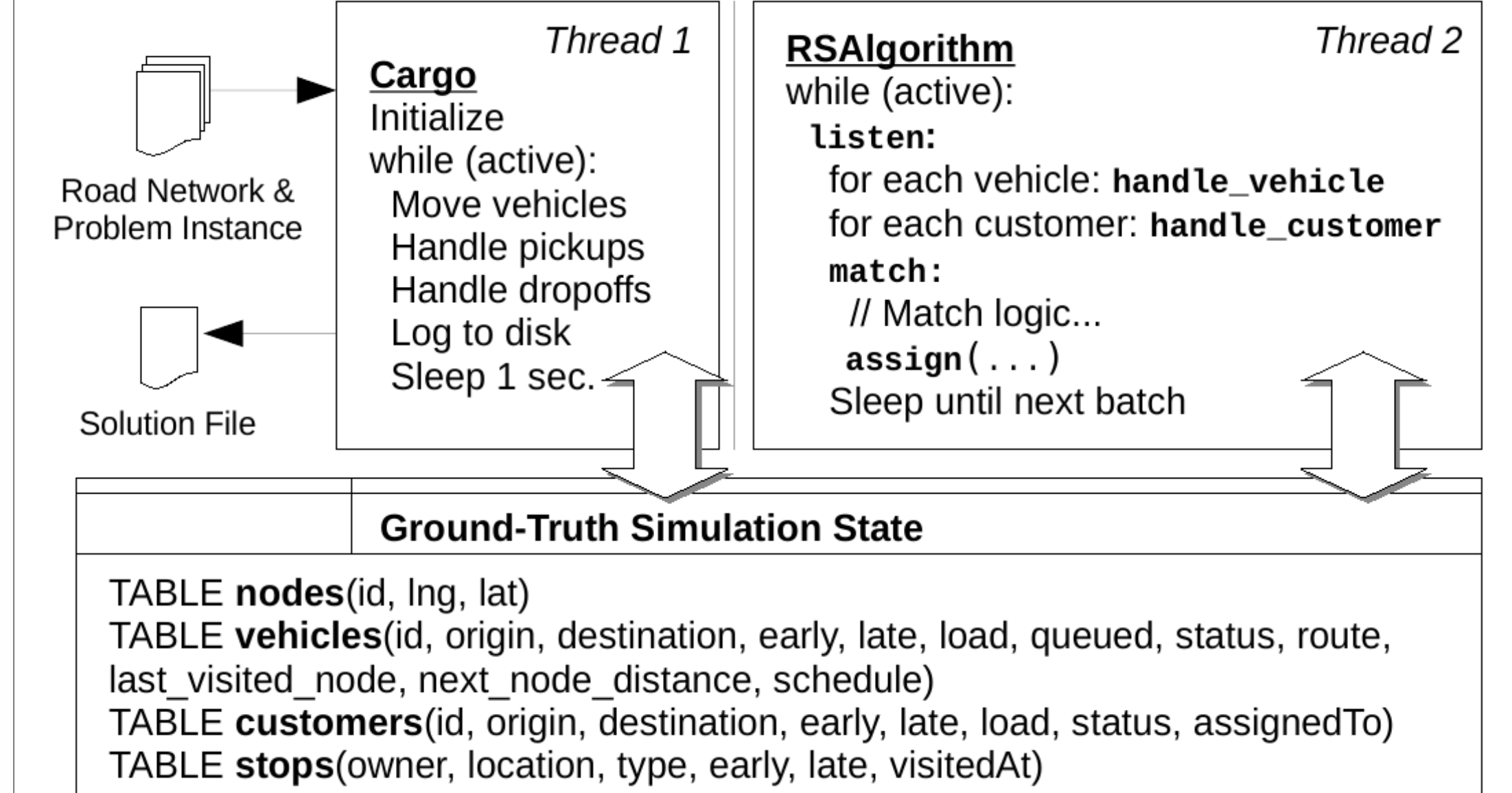
Descend through  $n$  randomized solutions, then select the best one. Descent procedures:

- Replace
- Swap
- Rearrange

Trip-Vehicle Grouping (TG)



## Our Ridesharing Simulator

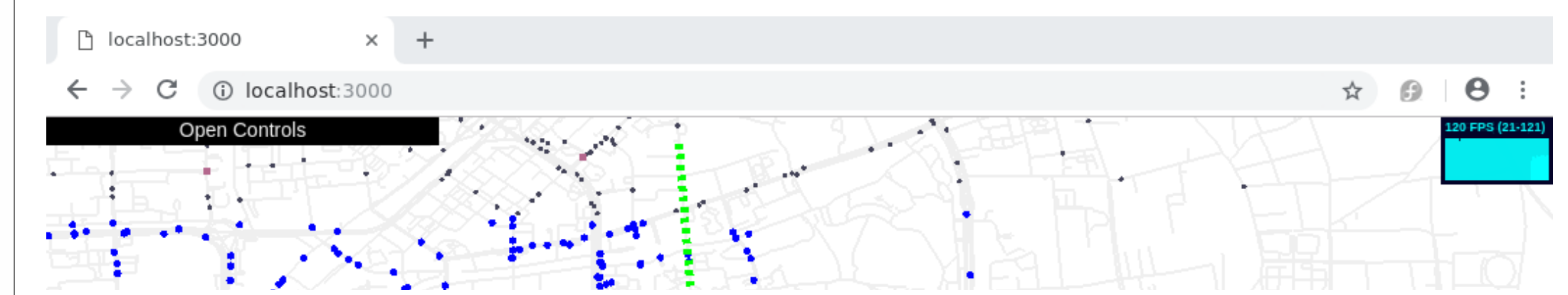


Modularized Ridesharing Algorithms:

```
class YourAlgorithm : public RSAAlgorithm {
public:
    YourAlgorithm();
    /* Your code in these methods */
    virtual void handle_customer(const Customer &);
    virtual void handle_vehicle(const Vehicle &);
    virtual void end();
    virtual void listen();
};
```

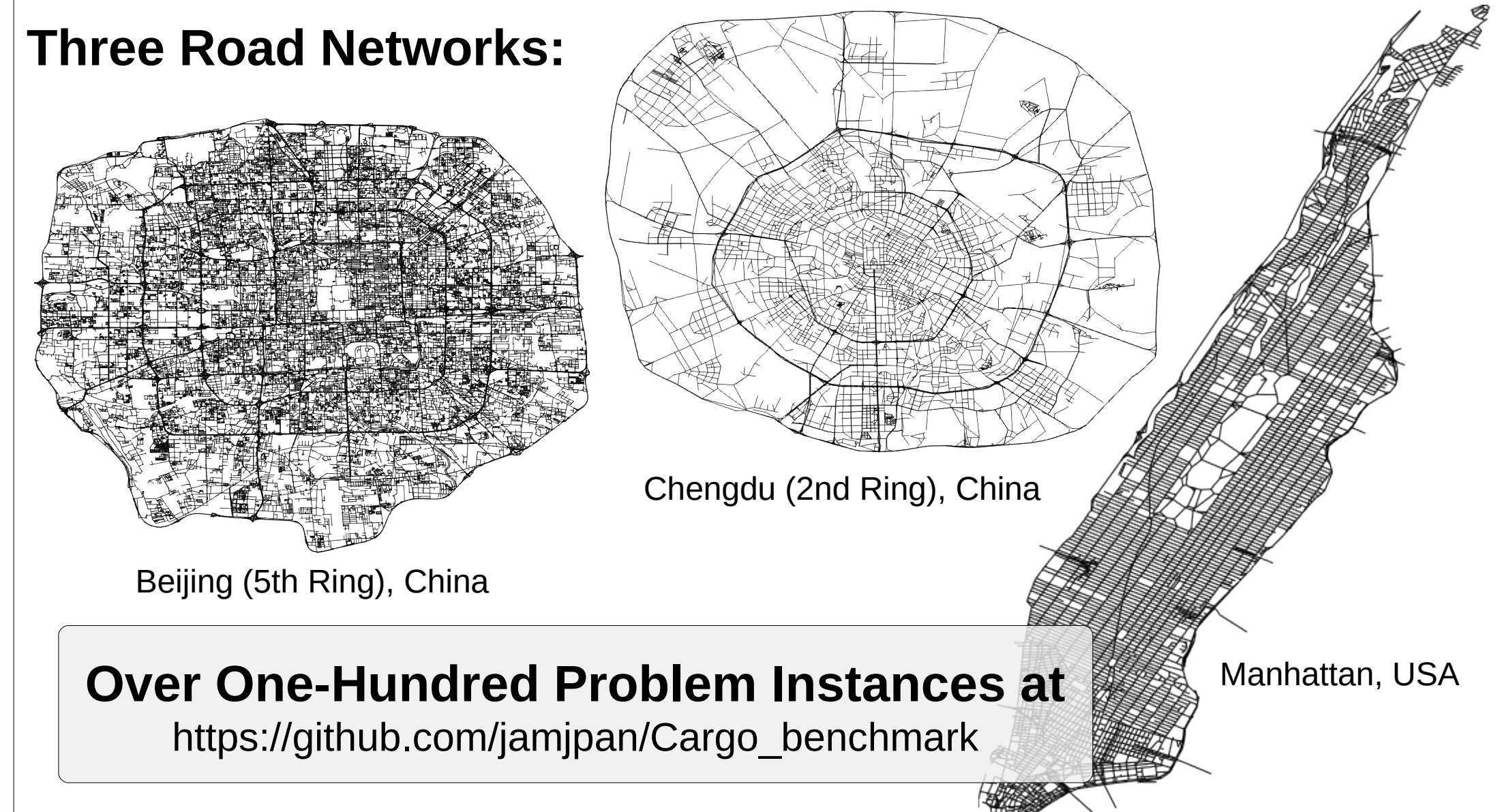
Visit <https://github.com/jamjpan/Cargo> for sourcecode and examples

WebGL Visualization with Live Updates:



## Our Benchmark Problem Instances

Three Road Networks:



Over One-Hundred Problem Instances at [https://github.com/jamjpan/Cargo\\_benchmark](https://github.com/jamjpan/Cargo_benchmark)